Matrix Representation of Groups in the Finite Fields $GF(2^n)$

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Abstract— The representation of mathematical fields can be accomplished by binary rows (or columns) of a binary triangular matrix as the Hamming's matrices, but this representation don't show the basic product properties of the fields, that is the nonzero elements of the fields forms a cyclic multiplicative group.

In this paper we show that the elements of the fields $GF(2^n)$, and their subgroups, can represent as square matrices by m – sequences, which satisfies the product properties as a cyclic group.

Index Term - Galois fields, m-sequences, cyclic groups, Orthogonal sequences.

I. INTRODUCTION

m- Linear Recurring Sequences

Let k be a positive integer and λ , λ_0 , λ_1 ,..., λ_{k-1} are

elements in the field F_q , then the sequence a_0, a_1, \dots is

called **non homogeneous linear recurring sequence of order** *kiff* :

$$a_{n+k} = \lambda_{k-1}a_{n+k-1} + \lambda_{k-2}a_{n+k-2} + \dots + \lambda_0 a_n + \lambda, \ \lambda_i \in F_q, \ i = 0, 1, \dots, k-1$$

or $a_{n+k} = \sum_{i=1}^{k-1} \lambda_i a_{n+i} + \lambda$ (1)

The elements $a_0, a_1, ..., a_{k-1}$ are called the **initial values** (or the vector $(a_0, a_1, ..., a_{k-1})$ is called **the initial vector**).

If $\lambda = 0$ then the sequence a_0, a_1, \dots

is called **homogeneous linear recurring sequence** (**H. L. R. S.**), except the zero initial vector, and the polynomial

$$f(x) = x^{k} + \lambda_{k-1} x^{k-1} + \dots + \lambda_{1} x + \lambda_{0}$$
(2)

is called the **characteristic polynomial.** In this study, we are limited to $\lambda_0 = 1$. [1]-[3]

II. THE IMPORTANCE OF THIS RESEARCH AND ITS OBJECTIVES

The elements of the fields $GF(2^n)$, and their subgroups, can be represented assquare matrices by m – sequences, which satisfies the multiplicative properties as a cyclic group, that is it will be useful in many other scientific branches. For example orthogonal sets in the forward and the inverse link of communications channels in the CDMA systems especially in the second (IS-95-CDMA), the third.... (CDMA200,...), the pilot channels, the Sync channels, and the Traffics channels in the present or in many other scientific branches in the future.

III. RESEARCH METHODS AND MATERIALS

Basic Definitions and Theorems

Definition 1.Let S be a nonempty set and a_0, a_1, \dots is

sequence from S and if there is r > 0 such that:

 $a_{n+r} = a_n$; $n \ge n_0$; $n_0 \ge 0$ (3)

Then this sequence is called Ultimately Periodic Sequence, and r is called a period of this sequence, the smallest positive integer between these r's is called the period of this sequence, and the smallest non negative n_0 such that:

$$a_{n+r} = a_n$$
; $n \ge n_0$; $n_0 \ge 0$

is called **Pre-Period**, [1][4] *Definition* **2.**The Ultimately Periodic Sequence

 a_0, a_1, \dots with the smallest period r is called a periodic

iff:
$$a_{n+r} = a_n$$
; $n = 0, 1, ... [1]-[4]$

Definition **3.**The complement of the binary vector

$$X = (x_1, x_2, ..., x_n) \text{ is the vector } \overline{X} = (x_1, x_2, ..., x_n),$$

when $\overline{x_i} = \begin{cases} 1 & \text{if } x_i = 0 \\ 0 & \text{if } x_i = 1 \end{cases}$. [1]-[4]

Definition 4. (Euler function φ). $\varphi(n)$ is the number of the natural numbers that are relatively prime with *n*.[5]-[8]

Definition 5. AnyPeriodic Sequence a_0, a_1, \dots over F_q

with prime characteristic polynomial is an orthogonal cyclic code and ideal auto correlation [1]-[10].

Definition 6. The binary periodic sequence $(a_i)_{i \in N}$, with the period *r* has the property of "Ideal Auto

Correlation" if and only if its periodic auto correlation $R_a(\tau)$ of the form:

$$R_a(\tau) = \begin{cases} r & ; \quad for \ \tau \equiv 0 \mod(r) \\ -1 & ; \quad otherwise \end{cases}$$

When:
$$R_a(\tau) = \sum_{t=0}^{r-1} (-1)^{a(\tau+t)+a(t)} [1], [2]$$

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Theorem1.

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i. If a_0, a_1, \dots is a homogeneous linear recurring sequence of order k in F_q , satisfies (1) then this sequence is periodic.

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ii. If this sequence is homogeneous linear recurring sequence, periodic with the period r, and its characteristic polynomial f(x) then r | ord f(x). [6]

iii. If the polynomial f(x) is primitive then the period of the sequence is $q^k - 1$, and this sequence is called m – sequence.

Lemma 2.(Fermat's theorem). If F is a finite field and has q elements then every elementa of F satisfies the

equation: $x^{q} = x$. [6],[9]

Theorem 3.For any primitive element p and any positive integer *n* there is a field *F*, which has p^n elements

and any two fields having $q = p^n$ Elements are isomorphic. [6],[9],[11] **Theorem 4.**

i. $(q^m - 1) | (q^n - 1) \Leftrightarrow m | n (4)$

ii. If F_q is a field of order $q = p^n$ then any subfield of

them of the order p^m and $m \mid n$ and by inverse if

 $m \mid n$ then in the field F_q there is a subfield of order

 p^m . [6],[9],[11]

Theorem 5. The number of irreducible polynomials in $F_q(x)$ of degree *m* and order *e* is $\varphi(e)/m$, if $e \ge 2$, when *m* is the order of *q* by mod *e*, and equal to 2 if m = a = 1, and equal to provide the provided of p = 1.

m = e = 1, and equal to zeroelsewhere. [6]-[9] **Theorem 6.** If g(x) is a characteristic prime polynomial of the (H. L. R. S.) a_0, a_1, \dots of degree k, and α is a root of g(x) in any splitting field of F_2 then the

general bound of the sequence is: $a_n = \sum_{i=1}^{k} C_i \left(\alpha^{2^{i-1}} \right)^n$.

[11],[12].

* The study here, is limited to the fields Galois $GF(2^n)$

IV. RESULTSAND DISCUSSION

A. First step

Theorem 7: Suppose a_0, a_1, \dots is a non zero homogeneous linear recurring sequence of order *k* over $F_2 = \{0,1\}$ and f(x) is its prime characteristic polynomial then the first $r = 2^k - 1$ bounds with all its cyclic shifts forming an additive group.

Proof: This sequence is periodic with period $r = 2^k - 1$. We suppose $\$ = \{S_1, S_2, ..., S_r\}$ where $S_1 = (a_1a_2...a_r)$

is the sequence of the first $r = 2^k - 1$ bounds, and $S_2 = (a_r a_1 \dots a_{r-1}), \dots, S_r = (a_2 a_3 \dots a_r a_1)$ are allits

cyclic shifts, and we suppose $O = S_0 = (0....0)$,

 $S = \$ \cup \{S_0\}$ and if α is a root of the prime polynomial f(x) and:

$$GF(2^{k}) = \left\{ \alpha^{i} : \alpha^{i} = \sum_{j=0}^{k-1} b_{j} \alpha^{j}, i = 0, 1, 2, \dots, 2^{k-1} \right\} \cup \{0\}, \ 0 = \sum_{i=0}^{k-1} 0 \alpha^{j}$$

And the function: $h: GF(2^k) \to S$ as following:

 $h(\alpha^{i}) = h(i) = h[b_0 \ b_1 \dots b_{k-1}] = \begin{bmatrix} b_0 \ b_1 & b_{k-1} \ b_k & b_{2^k-2} \end{bmatrix}$ Then h is one-to-one corresponding and:

$$\begin{cases} h(\alpha^{i} + \alpha^{j}) = h(\alpha^{i}) + h(\alpha^{j}) \\ h(m.\alpha^{i}) = m.h(\alpha^{i}), \quad m = 0 \text{ or } 1 \end{cases}$$

And *h* is Linear Transformation and isomorphism from the additive group $(GF(2^k),+)$ to the additive group (S,+).

B. Second Step

Theorem 8: Suppose a_1, a_2, \ldots is a non zerohomogeneous linear recurring sequence of order k in F_2 and f(x) is their primitive characteristic polynomial, S_1 is the initial bounds where $r = 2^k - 1$ and $\$ = \{S_1, S_2, \ldots, S_r\}$ are the all cyclic shifts. Let A is a matrix which its rows are elements \$ respectively, then $\{A^i, i = 1, \ldots, r\}$ is a cyclic multiplicative group, relatively to product of matrices, having the period of $S_1(x)$ and rows of A^i are the shifts to rows of A. **Proof:**

Suppose
$$A = \begin{bmatrix} S_1 \\ S_2 \\ \dots \\ S_r \end{bmatrix}$$
 and we will compute $A^2 = A \cdot A$.

Suppose the first row ω_1 in the matrix A^2 then:

$$\omega_1 = \sum_{i \in I} S_i$$

when *I* the set of all columns in *A* which does not start by zero, and we see that:

$$X,Y \in \$ \And X \neq Y \Longrightarrow X + Y \in \$$$

Then $\omega_l = S_l \in \$$

The second row ω_2 in A^2 is a result of shift *i* by 1 digit to the right, then: $\omega_2 = \sum_{i \in I} S_{i+1} = S_{l+1}$, and respectively we have $\omega_r = \sum_{i \in I} S_{i+r-1} = S_{l+r-1}$, when the

indexescomputed by *mod r*, then the rows of the matrix A^2 are shifts to rows of A, On other hand we suppose that $\$(x) = \{S_1(x), S_2(x), ..., S_r(x)\}$ then:



$$\omega_1(x) = \sum_{i \in I} S_i(x) ;$$

$$\omega_2(x) = \sum_{i \in I} S_{i+1}(x) = \sum_{i \in I} x S_i(x) ; \dots$$

$$\omega_r(x) = \sum_{i \in I} x^{r-1} S_i(x)$$

And:

$$\omega_{\mathbf{I}} = \sum_{i \in I} S_i(x) = \sum_{i \in I} x^{i-1} S_1(x) \implies \omega_{\mathbf{I}}(x) = S_1^2(x)$$

When: $S_1^2(x) \in \$(x)$, and the calculations are done by $\begin{pmatrix} & & \\ & & \\ & & \\ & & \\ & & \end{pmatrix}$

$$\left(\mod \left(x^{2^{-r_1}} - 1 \right) \right), \text{ And we have:}$$
$$\omega_2(x) = xS_1^2(x); \dots; \omega_r(x) = x^{r-1}S_1^2(x)$$

Suppose $[f_i(x)]$ denotes the row of coefficients of $f_i(x)$, respectively to increasing exponents of x, and which has the length r, then:

$$A = \begin{bmatrix} S_{1}(x) \\ S_{2}(x) \\ \vdots \\ S_{r}(x) \end{bmatrix} = \begin{bmatrix} S_{1}(x) \\ xS_{1}(x) \\ \vdots \\ x^{r-1}S_{1}(x) \end{bmatrix} ; A^{2} = \begin{bmatrix} S_{1}^{2}(x) \\ xS_{1}^{2}(x) \\ \vdots \\ x^{r-1}S_{1}^{2}(x) \end{bmatrix} ; \dots ;$$

$$A^{i} = \begin{bmatrix} S_{1}^{i}(x) \\ xS_{1}^{i}(x) \\ \vdots \\ x^{r-1}S_{1}^{i}(x) \end{bmatrix} , i = 1, 2, \dots, r$$

When: $S_1^i(x) \in \$(x); i = 1, ..., r$, then:

$$A = \begin{bmatrix} S_{1}(x) \\ S_{2}(x) \\ \vdots \\ S_{r}(x) \end{bmatrix} = \begin{bmatrix} 1 \\ x \\ \vdots \\ x^{r-1} \end{bmatrix} S_{1}(x) , \quad A^{2} = \begin{bmatrix} 1 \\ x \\ \vdots \\ x^{r-1} \end{bmatrix} S_{1}^{2}(x) , \dots$$

...,
$$A^{i} = \begin{bmatrix} 1 \\ x \\ . \\ x^{r-1} \end{bmatrix} S_{1}^{i}(x) , i = 1, 2, ..., r$$

Result 1: The period of the sequence A, A^2, A^3, \dots is equal to $ord(S_1(x))$ and divides $2^k - 1$.

Result 2: If $ord(S_1(x)) = 2^k - 1$ then S is representation to the field $GF(2^k)$.

Result 3: If $ord(S_1(x)) = l$ and $l | 2^k - 1$ then

 $l = 2^m - 1$ and $\langle A \rangle$ is a representation to one subgroup of order *l* in the field $GF(2^k)$ **Result 4:** If $ord(S_1(x)) + 1 = 2^l$ and $l \mid k$ then : $\langle A \rangle \cup \{O\}$ is representation to the Subfield $GF(2^l)$. **Result 5:** If $2^k - 1$ is prime then all elements of $S_i(x)$ are of the order $2^k - 1$ except only one of them that is of the order one

C. Third Step **Example 1:** If α is a root of the prime polynomial $f(x) = x^3 + x + 1$ and generates $GF(2^3)$ then the Binary representation of the elements of $GF(2^3)$ is:

$$\begin{array}{ll} \alpha & \to (1) = [010] \; ; \; \alpha^{2} = 1 + \alpha + \alpha^{2} \to (5) = [111] \\ \alpha^{2} & \to (2) = [001] \; ; \; \alpha^{6} = 1 + \alpha^{2} & \to (6) = [101] \\ \alpha^{3} = 1 + \alpha & \to (3) = [110] \; ; \; \alpha^{7} = 1 & \to (7) = [100] \\ \alpha^{4} = \alpha + \alpha^{2} \to (4) = [011] \; ; \; 0 & \to (0) = [000] \end{array}$$

Where (i), i = 0, 1, 2, ..., 7 is the symbol of the sequence *i*. The field $GF(2^3)$ contains two subfields : GF(2) and the same $GF(2^3)$ and the divisors of the number 7 are 1 and 7 consequently $GF(2^3)$ contains two multiplicative subgroups are: $GF^*(2)$ and $GF^*(2^3)$.

Suppose the Linear Recurring Sequence be $a_{n+3} = a_{n+1} + a_n$ or $a_{n+3} + a_{n+1} + a_n = 0$ (1)

Figure(1): Linear feedback register of degree3 generates sequence (1)

With the characteristic equation $x^3 + x + 1 = 0$ and the characteristic polynomial $f(x) = x^3 + x + 1$, which is a prime and generates F_{2^3} and if $\overline{x} = \alpha \in GF(2^3)$ is a root of f(x) then the solutions of characteristic equation are $\{\alpha^n, \alpha^{2n}, \alpha^{4n}\}$. The general solution of equation (1) isgiven by $a_n = \alpha^2 \cdot \alpha^n + \alpha^4 \cdot \alpha^{2n} + \alpha \cdot \alpha^{4n}$, and thesequence is periodic with the period $2^3 - 1 = 7$. For the initial position: $a_1 = 1, a_2 = 0, a_3 = 0$, then $S_1 = (1 \ 0 \ 0 \ 1 \ 0 \ 1 \ 1)$ and by the cyclic permutationson S_1 we have $\$ = \{S_1, S_2, S_3, S_4, S_5, S_6, S_7\}$ where:

The first three digits in each sequence are the initial position of the feedback register.



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Matrix Representation of Groups in the Finite Fields $GF(2^n)$

The matrix A of the cyclic permutations of S_1 is:

$$A = \begin{bmatrix} 1 & 0 & 0 & 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 & 0 & 1 & 0 \\ 0 & 1 & 1 & 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 \end{bmatrix}$$
 or briefly $A = A_{1} = \begin{bmatrix} 1 & 0 & 0 & 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 & 1 & 0 & 1 \\ - & - & - & - & - & - \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 \end{bmatrix}$

The first row in this matrix is called *The head row*.

If the function $h: GF(2^3) \rightarrow$ \$ where:

 $h(\alpha^{i}) = h(i) = h_{i}$ = [The row of the matrix A corresponding of the initial position I].

Then h_i is isomorphism from the group $(GF(2^3),+)$ on the group (\$,+).

I- In this matrix the head row is: $S_1 = \begin{bmatrix} 1 & 0 & 0 & 1 & 0 \end{bmatrix}$ and the head polynomial is:

 $h(1) = S_1(x) = 1 + x^3 + x^5 + x^6$.

We see that: $A_2 = A^2 = A$ and $\{O, A\}$ is a representation of the field GF(2), where O is zero matrix.

II- We suppose that;

$$B = B_1 = \begin{bmatrix} 1 & 1 & 0 & 0 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 & 0 & 1 & 0 \\ - & - & - & - & - & - \\ 1 & 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix}$$

In this matrix the head row is:

 $h(3) = S_2 = \begin{bmatrix} 1 & 1 & 0 & 0 & 1 & 0 & 1 \end{bmatrix}$ and the corresponding

head polynomial is:
$$S_2(x) = 1 + x + x^4 + x^6$$
, Thus
 $B_2 = B^2 = \begin{bmatrix} 1 & 1 & 1 & 0 & 0 & 1 & 0 \\ 0 & 1 & 1 & 1 & 0 & 0 & 1 \\ - & - & - & - & - & - \\ 1 & 1 & 0 & 0 & 1 & 0 & 1 \end{bmatrix}$

In this matrix the head row is:

 $h(5) = S_3 = \begin{bmatrix} 1 & 1 & 1 & 0 & 0 & 1 & 0 \end{bmatrix}$ and the corresponding head polynomial is: $S_2^2(x) = 1 + x + x^2 + x^5$, Thus

$$B_3 = B^3 = \begin{bmatrix} 0 & 1 & 1 & 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 1 & 1 & 0 & 0 \\ - & - & - & - & - & - \\ 1 & 1 & 1 & 0 & 0 & 1 & 0 \end{bmatrix}$$

In this matrix the head row is: $h(4) = S_4 = \begin{bmatrix} 0 & 1 & 1 & 1 & 0 & 0 & 1 \end{bmatrix}$ and the corresponding

head polynomial is: $S_2^3(x) = x + x^2 + x^3 + x^6$, Thus $B_4 = B^4 = \begin{bmatrix} 1 & 0 & 1 & 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 1 & 1 & 0 \\ - & - & - & - & - & - \\ 0 & 1 & 1 & 1 & 0 & 0 & 1 \end{bmatrix}$

In this matrix the head row is:

 $h(6) = S_5 = \begin{bmatrix} 1 & 0 & 1 & 1 & 1 & 0 & 0 \end{bmatrix}$ and the corresponding head polynomial is: $S_2^4(x) = 1 + x^2 + x^3 + x^4$, Thus

$$B_5 = B^5 = \begin{bmatrix} 0 & 1 & 0 & 1 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 \\ - & - & - & - & - & - \\ 1 & 0 & 1 & 1 & 1 & 0 & 0 \end{bmatrix}$$

In this matrix the head row is:

 $h(1) = S_6 = \begin{bmatrix} 0 & 1 & 0 & 1 & 1 & 1 & 0 \end{bmatrix}$ and the corresponding head polynomial is: $S_2^5(x) = x + x^3 + x^4 + x^5$, Thus $B_{6} = B^{6} = \begin{bmatrix} 0 & 0 & 1 & 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 & 0 & 1 & 1 \\ 1 & 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix}$

$$\begin{bmatrix} - & - & - & - & - & - \\ 0 & 1 & 0 & 1 & 1 & 1 & 0 \end{bmatrix}$$

In this matrix the head row is:

 $h(2) = S_7 = \begin{bmatrix} 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix}$ and the corresponding head polynomial is: $S_2^6(x) = x^2 + x^4 + x^5 + x^6$, Thus

$$B_7 = B^7 = \begin{bmatrix} 1 & 0 & 0 & 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 & 1 & 0 & 1 \\ - & - & - & - & - & - \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 \end{bmatrix}$$

In this matrix the head row is:

 $h(7) = S_1 = \begin{bmatrix} 1 & 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix}$ and the corresponding head polynomial is: $S_2^7(x) = 1 + x^3 + x^5 + x^6$, and

 $B_8 = B^8 = B$. We see that $\{O, B_1, B_2, B_3, B_4, B_5, B_6, B_7\}$ is a representation of $GF(2^3)$ and $ord(B_1) = ord(B_2) = ord(B_3) = ord(B_4) = ord(B_5) = ord(B_6) = 7$ The field F_{2^3} contains $\varphi(2^k - 1)/k = \varphi(2^3 - 1)/3 = 2$ third degree irreducible polynomials that are: $f(x) = x^3 + x + 1$ and its conjugate $g(x) = x^3 + x^2 + 1$, and g(x) is a prime polynomial then we can represent F_{2^3} by two different ways.

D. Fourth step

Example 2: If α is a root of the prime polynomial $f(x) = x^4 + x + 1$ and generates $GF(2^4)$ then the binary

representation of the elements of $GF(2^4)$ is:

 \rightarrow (1) = [0100]; $\alpha^9 = \alpha + \alpha^3$ α \rightarrow (5) = [0101] α^2 \rightarrow (2) = [0010]; $\alpha^{10} = 1 + \alpha + \alpha^2$ \rightarrow (10) = [1110] \rightarrow (3) = [000 1]; $\alpha^{11} = \alpha + \alpha^2 + \alpha^3 \rightarrow$ (11) = [0111] ~3 $\alpha^4 = 1 + \alpha \longrightarrow (4) = [1100]; \ \alpha^{12} = 1 + \alpha + \alpha^2 + \alpha^3 \rightarrow (12) = [1111]$ $\alpha^5 = \alpha + \alpha^2 \rightarrow (5) = [0\ 110] \ ; \ \alpha^{13} = 1 + \alpha^2 + \alpha^3 \rightarrow (13) = [1011]$ $\alpha^{6} = \alpha^{2} + \alpha^{3} \rightarrow (6) = [0 \ 0 \ 11] ; \alpha^{14} = 1 + \alpha^{3}$ \rightarrow (14) = [1 0 0 1] $\alpha^7 = 1 + \alpha + \alpha^3 \rightarrow (7) = [1101]$; $\alpha^{15} = 1$ \rightarrow (15) = [1000] $\alpha^8 = 1 + \alpha^2 \qquad \rightarrow (8) = [1010]; \quad 0$ \rightarrow (0)=[0000]

Where (i), i = 0, 1, 2, ..., 15 is the symbol of the sequence i. The field $GF(2^4)$ contains three subfields are: GF(2),



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 $h(\alpha^{i}) = h(i) = h_{i}$ = [The row of the matrix A corresponding

 $GF(2^2)$ and the same $GF(2^4)$ and the divisors of the number 15 are 1, 3, 5 and 15 consequently $GF(2^4)$ contains four multiplicative subgroups : $GF^*(2)$, $GF^*(2^2)$, one group of the order 5 and $GF^*(2^4)$. Suppose the Linear Recurring Sequence:

 $a_{n+4} = a_{n+1} + a_n \text{ or } a_{n+4} + a_{n+1} + a_n = 0$ (2)



Figure(2): Linear feedback register of degree 4 generates sequence (2)

With the characteristic equation $x^4 + x + 1 = 0$ and the characteristic polynomial $f(x) = x^4 + x + 1$, which is prime and generates F_{2^4} and if $\overline{x} = \alpha \in GF(2^4)$ is a root of f(x) then the solutions of characteristic equation are $\{\alpha, \alpha^2, \alpha^4, \alpha^8 = \alpha^2 + 1\}$ and the general solution of equation (2) is given by:

$$a_n = c_1 \alpha^n + c_2 a^{2n} + c_3 \alpha^{4n} + c_4 \alpha^{8n}$$

For $a_1 = 1$, $a_2 = 0$, $a_3 = 1$, $a_4 = 0$ we have: $a_n = \alpha^9 \alpha^n + \alpha^{12} a^{2n} + \alpha^9 \alpha^{4n} + \alpha^3 \alpha^{8n}$

and the sequence is periodic with the period $2^4 - 1 = 15$. $s_1 = (1,0,1,0,1,1,1,1,0,0,0,1,0,0,1); s_2 = (1,1,0,1,0,1,1,1,1,0,0,0,1,0,0);...$

$$\dots; s_{15} = (0,1,0,1,1,1,1,0,0,0,1,0,0,1,1)$$

The first four digits in each sequence are the initial position of the feedback register.

The matrix of the cyclic permutations of S_1 is the following matrix:

	[1	0	1	0	1	1	1	1	0	0	0	1	0	0	1	
	1	1	0	1	0	1	1	1	1	0	0	0	1	0	0	
	0	1	1	0	1	0	1	1	1	1	0	0	0	1	0	
	0	0	1	1	0	1	0	1	1	1	1	0	0	0	1	
	1	0	0	1	1	0	1	0	1	1	1	1	0	0	0	
	0	1	0	0	1	1	0	1	0	1	1	1	1	0	0	
	0	0	1	0	0	1	1	0	1	0	1	1	1	1	0	
A =	0	0	0	1	0	0	1	1	0	1	0	1	1	1	1	
	1	0	0	0	1	0	0	1	1	0	1	0	1	1	1	
	1	1	0	0	0	1	0	0	1	1	0	1	0	1	1	
	1	1	1	0	0	0	1	0	0	1	1	0	1	0	1	
	1	1	1	1	0	0	0	1	0	0	1	1	0	1	0	
	0	1	1	1	1	0	0	0	1	0	0	1	1	0	1	
	1	0	1	1	1	1	0	0	0	1	0	0	1	1	0	
	0	1	0	1	1	1	1	0	0	0	1	0	0	1	1	I
Or briefly	/:															
•	Γ	1 ()	1 (0	1	1	1	1	0	0	0	1	0	0	1
4 - 4 -	_	1	1	0	1	0	1	1	1	1	0	0	0	1	0	0
$A = A_1$	- -				_	_	_	_	_	_	_	_	_	_	_	_
)	1	0	1	1	1	1	0	0	0	1	0	0	1	1
The first	ros	w ir	h th	ie n	nat	riv	ic	11	۵đ	ть	o h	oad	l ro	147		

The first row in this matrix is called *The head row*. If the function $h: GF(2^4) \rightarrow \$$ where:

of the initial position *i*], then h_i is isomorphism from the additive group $(GF(2^4),+)$ on the additive group (\$,+). In this matrix the head row is: And the corresponding head polynomial is: $S_1(x) = 1 + x^2 + x^4 + x^5 + x^6 + x^7 + x^{11} + x^{14}$, Thus: In this matrix the head row is: and the corresponding head polynomial is: $S^{3}(x) = 1 + x + x^{3} + x^{5} + x^{6} + x^{7} + x^{8} + x^{12}$ Thus: In this matrix the head row is: and the corresponding head polynomial is: $S^{3}_{\cdot}(x) = 1 + x + x^{3} + x^{5} + x^{6} + x^{7} + x^{8} + x^{12}$ Thus: In this matrix the head row is: and the corresponding head polynomial is: $S^{4}(x) = 1 + x + x^{5} + x^{8} + x^{9} + x^{11} + x^{13} + x^{14}$ Thus: In this matrix the head row is: and the corresponding head polynomial In this matrix the head row is: and the corresponding head polynomial

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Matrix Representation of Groups in the Finite Fields $GF(2^n)$

In this matrix the head row is:

is:
$$S_{1}^{13}(x) = x^2 + x^5 + x^6 + x^8 + x^{10} + x^{11} + x^{12} + x^{13}$$
 Thus:

In this matrix the head row is: and the corresponding head polynomial is: $S^{14}(x) = x + x^3 + x^4 + x^5 + x^6 + x^{10} + x^{13} + x^{14}$ Thus: In this matrix the head row is: and the corresponding head polynomial is: $S_{1}^{15}(x) = x^3 + x^6 + x^7 + x^9 + x^{11} + x^{12} + x^{13} + x^{14}$ Thu s: $A_{16} = A^{16} = A$. By considering $h_0 = O$ we see that: $ord(A_1) = ord(A_2) = ord(A_4) = ord(A_7) = ord(A_8) =$ $= ord(A_{11}) = ord(A_{13}) = ord(A_{14}) = 15$ According to above calculations: The set: $\mathbf{A} = \{O, A, A_2, A_3, \dots, A_{15}\}$ is an additive group (Table 1) and the set $\mathbf{A}^* = \{A, A_2, A_3, \dots, A_{15}\}$ is a multiplicative group of order 15 the identity is A_{15} (Table 2). Then the set **A** is a representation of $GF(2^4)$. The field F_{2^4} contains $\varphi(2^k - 1)/k = \varphi(2^4 - 1)/4 = 2$ third degree irreducible polynomials that are $f(x) = x^4 + x + 1$ and its conjugate $g(x) = x^4 + x^3 + 1$, also g(x) is a prime polynomial then we can represent F_{2^4} by two different ways.

is: $S_2^2(x) = 1 + x + x^2 + x^6 + x^9 + x^{10} + x^{12} + x^{14}$ Thus:



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In this matrix the head row is:

is:
$$S_2^3(x) = 1 + x + x^3 + x^5 + x^6 + x^7 + x^8 + x^{12}$$
 Thus:

In this matrix the head row is:

is:
$$S_{2}^{5}(x) = x^{3} + x^{6} + x^{7} + x^{9} + x^{11} + x^{12} + x^{13} + x^{14}$$

Thus: $B_6 = B^6 = B$.

Then the set $\mathbf{B} = \{O, B_1, B_2, B_3, B_4, B_5\}$ is not an additive group because $B + B_2$ does not belongs to the set and: $ord(B_1) = ord(B_2) = ord(B_3) = ord(B_4) = 5$. Also, $\mathbf{B}^* = \{B_1, B_2, B_3, B_4, B_5\}$ is a multiplicative group of order 5 and the identity in their is $B_5 = A_{1,5}$. (Table 3).

The field F_{2^4} contains $\varphi(2^k - 1)/k = \varphi(2^4 - 1)/4 = 2$ third degree irreducible polynomials that are:

 $f(x) = x^4 + x + 1$ and its conjugate $g(x) = x^4 + x^3 + 1$. Also, g(x) is a prime polynomial then we can represent the group $\{B, B_2, B_3, B_4, B_5\}$ by two different ways.

III- Suppose the matrix:

11																
	۲O	1	1	0	1	0	1	1	1	1	0	0	0	1	0	
C = C =	0	0	1	1	0	1	0	1	1	1	1	0	0	0	1	
$c - c_1 - c_1$	-	_	_	_	_	_	_	_	_	_	_	_	_	_	-	
	1	1	0	1	0	1	1	1	1	0	0	0	1	0	0	
In this m	atriz	k th	e he	ead	rov	v is	:									
$h(5) = S_3$	3 =	0	1	1 (C	1 (0 3	1 1	1	1	0	0	0	1	0]	
And the	corr	esp	ond	ing	he	ad j	poly	nor	nia	l is:						
$S_3(x) =$	<i>x</i> +	$-x^2$	+.	x^4	+x	;6 +	$+x^7$	′+.	x^8	+x	9+	x^1	³ Tł	nus		
	[0	1	1	1	1	0	0	0	1	0	0	1	1	0	1]	
$C_{1} = C^{2}$	1	0	1	1	1	1	0	0	0	1	0	0	1	1	0	
$c_2 - c_{-}$		_	_	_	_	_	_	_	_	_	_	_	_	-	-	
	1	1	1	1	0	0	0	1	0	0	1	1	0	1	0	
In this m	natri	x th	ne h	ead	ro	w is	s:									
h(11) = 3	S ₁₃ :	=[0	1	1	1	1	0	0	0	1 () () 1	1	0	1]	

and the corresponding head polynomial is:

And the set $\{O, C_3\}$ represent the field $\{O, C_3\}$. Thus we can represent it by two different ways.

V. RESULTS AND RECOMMENDATIONS

- 1. The fields $GF(2^n)$ and their subfield can be represented by square matrices.
- 2. The multiplicative group $GF^*(2^n)$ and their subgroups can be represented by square matrices.
- 3. The equations of the degree less than or equal to n on GF(2) can also be solved by square matrices.
- 4. Building encoders on the field F_q when $q \mid n$ are recommended for further study.

APPENDIX



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+	C	C_2	C_3	•	C	C_2	C_3
C	0	C_3	C_2	\mathcal{C}	C_2	C_3	\mathcal{C}
C_2	C_3	0	C	\mathcal{C}_2	C_3	Ĉ	C_2
C_3	\mathcal{C}_2	C	0	C_3	C	C_2	C_3

 Table 1: Addition on the additive group A

•	Α	A_2	A_3	A4	A_5	A_6	A7
А	A_2	A_3	A_4	A_5	A_6	A_7	A_8
A_2	A_3	A4	A_5	A_6	A7	A_8	Ag
A_3	A4	A_5	A_6	A7	A ₈	Ag	A_{10}
A_4	A_5	A_6	A_7	A_8	A9	A_{10}	A_{11}
A_5	A_6	A_7	A_8	A_9	A_{10}	A_{11}	A_{12}
A_6	A_7	A_8	A_9	A_{10}	A_{11}	A_{12}	A_{13}
A_7	A_8	A_9	A_{10}	A_{11}	A_{12}	A_{13}	A ₁₄
A_8	A_9	A_{10}	A_{11}	A_{12}	$A_{\rm B}$	A_{14}	A_{15}
A9	A_{10}	A_{11}	A_{12}	A_{13}	A_{14}	A_{15}	Α
A_{10}	A_{11}	A_{12}	A_{13}	A_{14}	A_{15}	Α	A_2
A_{11}	A_{12}	A_{13}	A_{14}	A_{15}	A	A_2	A_3
A_{12}	A ₁₃	A ₁₄	A_{15}	A	A_2	A_3	A4
$A_{\rm B}$	A_{14}	A_{15}	A	A_2	A_3	A_4	A_5
A ₁₄	A ₁₅	A	A_2	A_3	<i>A</i> ₄	A_5	A_6
A_{15}	A	A_2	A_3	A_4	A_{5}	A_6	A_7

Table 2: Multiplication on the multiplicative group A^{*}

	A_8	<i>A</i> 9	A_{10}	<i>A</i> ₁₁	A_{12}	$A_{\rm B}$	A ₁₄	A ₁₅
	Ag	A_{10}	A_{11}	A_{12}	A_{13}	A_{14}	A_{15}	Α
	A_{10}	A_{11}	A_{12}	A_{13}	A_{14}	$A_{\rm LS}$	Α	A_2
	A_{11}	A_{12}	A ₁₃	A ₁₄	A_{15}	Α	A_2	<i>A</i> ₃
	A_{12}	$A_{\rm B}$	A_{14}	A_{15}	Α	A_2	A_3	A4
ľ	A ₁₃	A ₁₄	A ₁₅	Α	A_2	A_3	A_4	A_5
Γ	A_{14}	A_{13}	Α	A_2	A_3	A_4	A_5	A_6
Ī	A_{15}	A	A_2	A_3	A_4	A_5	A_6	A_7
	Α	A_2	A_3	A_4	A_5	A_6	A_7	A ₈
	A_2	A_3	A_4	A_5	A_6	A_7	A_8	A9
	A_3	A_4	A_5	A_6	A_7	A_8	A_9	A_{10}
	A_4	A_5	A_6	A_7	A_8	A9	A_{10}	A_{11}
	A_5	A_6	A_7	A_8	A_9	A_{10}	A_{11}	A_{12}
	A_6	A_7	A_8	A_9	A_{10}	A_{11}	A_{12}	A ₁₃
	A_7	A_8	A_9	A_{10}	A_{11}	A_{12}	A ₁₃	A ₁₄
	A_8	A9	A_{10}	A_{11}	A_{12}	$A_{\rm B}$	A_{14}	A ₁₅

 Table 3:Multiplication on multiplicative group B^{*}

Tables 4: Addition on C and multiplication on C*

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